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CLAIMS: The following is a listing of all claims in the application with their status and the text for all active claims.

- 1. (CANCELED)
- 2. (CANCELED)
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- 25. (CANCELED)
- 26. (CANCELED) 27. (CANCELED)
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- 28. (CANCELED)
- 29. (CANCELED)
- 30. (CANCELED)
- 31. (CANCELED)
- 32. (CANCELED)
- 33. (CURRENTLY AMENDED) A computer-implemented method for improving compression for storage of a plurality of parallel data element sequences comprising:

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- (a) creating a dictionary of unique values for each of said data element sequences, wherein each dictionary associates a numeric index with each unique value,
- (b) forming an n-ary tree with leaf and interior nodes wherein:
  - (1) each said leaf node corresponds to one of said dictionaries,
  - (2) each said interior node associates a numeric index with tuples of numeric indexes from other subordinate leaf or interior nodes,
  - (3) one or more interior nodes are capable of storing store one or more sequences of mutually-consecutive tuples by representing said sequences in a form that uses less storage space than representing said sequences as individual tuples, and
  - (4) one or more interior nodes are capable of perform at least one of:
    - recording the addition of a tuple that extends a tuple sequence by modifying one
      or more fields in the representation of said sequence that are capable of
      representing the length of said sequence, or
    - ii. recording the addition of a tuple that invalidates an existing tuple sequence by splitting said tuple sequence into one or more subsequences, wherein none of the tuples of said subsequences contain any element of said added tuple, or
    - iii. recording the addition of a tuple that has not been previously added to said interior node, wherein said added tuple does not extend a tuple sequence, by adding said tuple to a tuple collection,-or
    - iv. any combination of two or more of i, ii, and iii,

wherein the forming step comprises:

- (c) defining a problem space comprising:
  - (1) a set of states such that each said state contains an n-ary tree design, comprising a set of leaves and zero or more interior nodes, each said interior node with zero or more other nodes as children, and
  - (2) a value function, giving a numeric ranking of the value of any state's <u>n-ary tree</u> design, wherein said numeric ranking is used to either.
    - a determine if said state's n-ary tree design is acceptable, or
    - b rank the n-ary tree designs of states that can be reached by applying operators to a state, to select an operator,
- (d) defining one or more operators that transform one state to another, and
- (e) searching the <u>said problem</u> space, starting from an initial state and applying operators to move to other states until a state with an acceptable n-ary tree design is reached.
- 34. (CURRENTLY AMENDED) The method of claim 3353, wherein said method for arranging an n ary tree uses an estimate of interior node size, from uses a function of the sizes of said interior node's child nodes.

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- 35. (PREVIOUSLY ADDED) The method of claim 33, wherein said tree-arrangement method uses an operator that joins two or more nodes (leaf nodes or interior nodes) under an interior node
- 36. (PREVIOUSLY ADDED) The method of claim 33, wherein each unique value of a leaf node or each unique tuple of an interior node is capable of being associated with a count of the number of times that value or tuple of values occurred in the parallel data element sequences.
- 37. (PREVIOUSLY ADDED) The method of claim 33, wherein said method is used to compress a table, wherein said parallel data element sequences represent fields of said table's records.
- 38. (CANCELLED) The method of claim 33, wherein said value function gives a value which indicates whether a state's n-ary tree design is acceptable.
- (CURRENTLY AMENDED) A computer-implemented method for improving compression for storage of a plurality of parallel data element sequences comprising:
  - (a) creating a dictionary of unique values for each of said data element sequences, wherein each dictionary associates a numeric index with each unique value,
  - (b) forming one or more n-ary trees with leaf and interior nodes wherein:
    - at least one of said leaf nodes is distinct from, and represents a subset of values from one of said dictionaries,
    - (2) each interior node associates a numeric index with tuples of numeric indexes from other subordinate leaf or interior nodes,

wherein the forming step comprises:

- (c) defining a problem space comprising:
  - (1) a set of states such that each said state contains an n-ary tree design, comprising a set of leaves and zero or more interior nodes, each <u>said interior node</u> with zero or more other nodes as children, and
  - (2) a value function, giving a numeric ranking of the value of any state's <u>n-ary tree</u> design, wherein said numeric ranking is used to either.
    - a. determine if said state's n-ary tree design is acceptable, or
    - b. rank the n-ary tree designs of states that can be reached by applying operators to a state, to select an operator,
- (d) defining one or more operators that transform one state to another, and
- (e) searching the <u>said problem</u> space, starting from an initial state and applying operators to move to other states until a state with an acceptable n-ary tree design is reached.
- 40. (PREVIOUSLY ADDED) The method of claim 39, wherein said method for arranging an n-ary tree uses an estimate of interior node size, from a function of the sizes of said interior node's child nodes.

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- 41. (PREVIOUSLY ADDED) The method of claim 39, wherein said tree-arrangement method uses an operator that joins two or more nodes (leaf nodes or interior nodes) under an interior node
- 42. (PREVIOUSLY ADDED) The method of claim 39, wherein each unique value of a leaf node or each unique tuple of an interior node is capable of being associated with a count of the number of times that value or tuple of values occurred in the parallel data element sequences.
- 43. (CURRENTLY AMENDED) A computer-implemented method for storage of a plurality of parallel data element sequences, and efficiently processing elements from a subset of said sequences, comprising:
  - (a) creating a dictionary of unique values for each of said data element sequences, wherein each dictionary associates a numeric index with each unique value,
  - (b) forming one or more n-ary trees with leaf and interior nodes wherein:
    - (1) each leaf node corresponds to one of said dictionaries,
    - (2) each interior node associates a numeric index with tuples of numeric indexes from other subordinate leaf or interior nodes.
    - (3) a gate field is defined for one or more interior nodes,
  - (c) processing the leaves corresponding to said subset of sequences by:
    - setting the value of said gate field for each said interior node, to indicate which of said interior node's branches lead to leaf nodes in said subset,
    - (2) following paths that lead to said leaf nodes, and
    - (3) processing said elements in said leaf nodes encountered,

wherein the forming step comprises:

- (d) defining a problem space comprising:
  - (1) a set of states such that each said state contains an n-ary tree design, comprising a set of leaves and zero or more interior nodes, each <u>said interior node</u> with zero or more other nodes as children, and
  - (2) a value function, giving a numeric ranking of the value of any state's <u>n-ary tree</u> design, wherein said numeric ranking is used to either:
    - a. determine if said state's n-ary tree design is acceptable, or
    - tank the n-ary tree designs of states that can be reached by applying operators to a state, to select an operator.
- (e) defining one or more operators that transform one state to another, and
- (f) searching the <u>said problem</u> space, starting from an initial state and applying operators to move to other states until a state with an acceptable n-ary tree design is reached.

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- 44. (PREVIOUSLY ADDED) The method of claim 43, wherein said method for arranging an n-ary tree uses an estimate of interior node size, from a function of the sizes of said interior node's child nodes.
- 45. (PREVIOUSLY ADDED) The method of claim 43, wherein said tree-arrangement method uses an operator that joins two or more nodes (leaf nodes or interior nodes) under an interior node.
- 46. (PREVIOUSLY ADDED) The method of claim 43, wherein each unique value of a leaf node or each unique tuple of an interior node is capable of being associated with a count of the number of times that value or tuple of values occurred in the parallel data element sequences.
- 47. (PREVIOUSLY ADDED) The method of claim 43, where said processing includes using values or tokens at said leaf nodes to reconstruct a subset of a stored record.
- 48. (PREVIOUSLY ADDED) The method of claim 47, further including the step of adding one or more of said reconstructed record subsets to another tree.
- 49. (PREVIOUSLY ADDED) The method of claim 33 wherein at least one of said leaf nodes is distinct from, and represents a subset of values from one of said dictionaries.
- 50. (PREVIOUSLY ADDED) The method of claim 33, further including a method for efficiently processing elements from a subset of said parallel data element sequences, comprising:
  - (a) defining a gate field for one or more interior nodes,
  - (b) processing the leaves corresponding to said subset of sequences by:
    - (1) setting the value of said gate field for each said interior node, to indicate which of said interior node's branches lead to leaf nodes in said subset,
    - (2) following paths that lead to said leaf nodes, and
    - (3) processing said elements in said leaf nodes encountered.
- 51. (PREVIOUSLY ADDED) The method of claim 39, further including a method for efficiently processing elements from a subset of said parallel data element sequences, comprising:
  - (a) defining a gate field for one or more interior nodes,
  - (b) processing the leaves corresponding to said subset of sequences by:
    - setting the value of said gate field for each said interior node, to indicate which of said interior node's branches lead to leaf nodes in said subset,
    - (2) following paths that lead to said leaf nodes, and
    - (3) processing said elements in said leaf nodes encountered.
- 52. (PREVIOUSLY ADDED) The method of claim 49, further including a method for efficiently processing elements from a subset of said parallel data element sequences, comprising:
  - (a) defining a gate field for one or more interior nodes,
  - (b) processing the leaves corresponding to said subset of sequences by:

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- (1) setting the value of said gate field for each said interior node, to indicate which of said interior node's branches lead to leaf nodes in said subset,
- (2) following paths that lead to said leaf nodes, and
- (3) processing said elements in said leaf nodes encountered.
- 53. (NEW) The method of claim 33, wherein said problem space is represented by a state containing leaves and zero or more interior nodes, wherein said state representation can be modified by operators to represent different states, and said value function is a function of the actual or estimated sizes of a state's interior and/or leaf nodes.